Mantis 1356
P1800-2012

Motivation

This Mantis item enables the use of Java style interfaces in the place of true multiple inheritance (MI) as implemented in C++. Please see Dave Rich’s paper titled “The Problems with Lack of Multiple Inheritance in SystemVerilog and a Solution” for a good history and need for interfaces. We have chosen the Java approach, with some subtle variations needed for SVTB, because of the integration complexity associated with a full MI solution and because we see the Java solution as meeting the needs of an MI approach for SystemVerilog. The restrictions that we have chosen basically limit interface classes to classes with pure virtual methods. WRT the diamond name resolution issue highlighted in Dave’s paper, we choose to “hide”, or in other words not inherit, parameters and other name scoped tokens of the interface class. These types can still be accessed with the class scope operator ‘::’, they are just not inherited. We choose to introduce the keyword ‘interface’ and the concept of ‘interface classes” rather than Dave’s suggested ‘virtual <classname>” as this best represents the intent of this new functionality. We do not believe this will conflict with SV interfaces or overuse that keyword as this new functionality will be introduced and discussed in the context of being an ‘interface class’.

(NOTE: There were several mantis issues opened up to deal with resolving other aspects of the LRM that need attention before we can bottom out on Interface Class refinement. In most of these cases, this will lead to a restriction to specific Interface Class features until we can resolve them. Where pertinent, I have noted the issue, the restriction, and the open mantis ticket).

(NOTE: BNF will be written once the details of this spec are near completion)

8.1 General

This clause describes the following:
— Class definitions
— Virtual classes and methods
— Polymorphism
— Parameterized classes
— Interface Classes

8.3 Syntax

class_declaration ::= // from A.1.2
[ virtual ] class [ lifetime ] class_identifier [ parameter_port_list ]
[ extends class_type [ { list_of_arguments } ] ]
[ implements interface_class_type [ { list_of_arguments } ] {, interface_class_type [ { list_of_arguments } ] } ]
[ { class_item } :
endclass [ : class_identifier]
interface_class_type: class_identifier[parameter_value_assignment]
…….

Add sub-clause 8.25 (and increment 8.25→8.27 by one number) to the Classes Clause as follows

8.25 Interface Classes

SystemVerilog introduces a type of class called an interface class. This is not to be confused with the interface construct (see section 3.5 Interfaces). An interface class can be thought of as a prototype or
skeleton of a class where the methods within it are outlined, but not defined, thereby setting up a framework of how the class should be implemented. An interface class shall contain only methods of pure virtual type (see 8.20). Other unrelated classes can implement the interface class through the implements keyword but must fully define the methods with the exception of virtual classes. A virtual class that is implementing an interface class is not required to fully define the methods, however they must be fully defined by descendent classes.

A class may be declared to directly implements one or more interface classes, meaning that any instance of the class implements all the pure virtual methods specified by one or more interface classes. A class must implement all of the pure virtual methods that are prototyped by these interface classes. This (multiple) interface inheritance allows classes to support (multiple) common behaviors without sharing any implementation.

An interface class makes it unnecessary for related classes to share a common abstract superclass or for that superclass to contain all method definitions needed by child classes. An interface class may be declared to be an extension of one or more other interface classes, meaning that it implicitly specifies all the member types, pure virtual methods and constants of the interface classes it extends, except for any member types and constants that it may hide.

A variable whose declared type is an interface class type may have as its value a reference to any instance of a class which implements the specified interface class (see 8.21 Polymorphism). It is not sufficient that the class happens to implement all the pure virtual methods of the interface class; the class or one of its superclasses must actually be declared to implement the interface class through the implements keyword, or else the class is not considered to implement the interface class.

The following is a simple example of interface classes.

```verbatim
interface class PutImp#(type T = logic, string MyPutSignal = "a");
    pure virtual task void put(T MyPutSignal);
endclass

interface class GetImp#(type T = logic);
    pure virtual task T get();
endclass

class Fifo#(type T = logic, DEPTH = 1) implements PutImp#(T), GetImp#(T);
    T [DEPTH-1:0] myFifo;
    virtual task void put([T] a);
    // Put implementation
    virtual task T get();
    // Get implementation
endclass
```
There is a difference between how classes, virtual classes, and interface classes inherit each other. The following highlights these differences:

- An interface class
  - may extend zero or more interface classes;
  - may not implement an interface class
  - may not extend a non-interface class.

- A class or virtual class
  - may extend at most one other class or virtual class;
  - may implement zero or more interface classes;
  - may not extend an interface class.
  - may both extend a class and implement interface classes.

Conceptually an `extends` is considered a means to extend the content of the parent class while an `implements` is considered a contract on behalf of the implementing class to supply the definition of the interface class. Whenever the keyword `implements` is used the implementing class shall supply a definition or an error will be issued. An interface class may be extended to an interface class, meaning that the sub interface class can have additional methods outlined but may not define any of them. A virtual class can extend one class and/or implement one or more interface classes. Because virtual classes are abstract they may or may not choose to fully define the methods from their parent class. Therefore virtual classes may provide a means to create partial implementations of classes (See 8.25.5 Partial implementations). A class can only implement interface classes. It shall be an error to extend a class from an interface class.

The following example a class is both extending a base class and implementing two interface classes:

```verilog
class MyQueue (type T = logic, DEPTH = 1, string MyPutSignal = "a");
  T[DEPTH-1:0] PipeQueue[$];
  virtual function void deleteQ();
  // Delete implementation
endfunction
class Fifo extends MyQueue#(T, DEPTH), implements PutImp#(T), GetImp#(T);
  virtual task void put(T b);
  // Put implementation
endtask
  virtual task T get();
  // Get implementation
endtask
class MyQueue (type T = logic, DEPTH = 1, string MyPutSignal = "a");
  T[DEPTH-1:0] PipeQueue[$];
  virtual function void deleteQ();
  // Delete implementation
endfunction
class Fifo extends MyQueue#(T, DEPTH), implements PutImp#(T), GetImp#(T);
  virtual task void put(T b);
  // Put implementation
endtask
  virtual task T get();
  // Get implementation
endtask
endclass
```

In this example, the `Fifo` child class is extending the `MyQueue` base class which has a parameterized queue and an associated `deleteQ()` method. This property and method are inherited in the `Fifo` class.
addition the Fifo class is also implementing the PutImp and GetImp interface classes and defining the put and get methods respectively.

In the following example we demonstrate that multiple types can be parameterized in the class definition and the resolved types used in the implemented classes PutImp and GetImp.

class Fifo#(type T_in = logic, type T_out = logic, DEPTH = 1) implements PutImp#(T_in), GetImp#(T_out);
  T_in [DEPTH-1:0] myFifo;
  virtual task void put(T_in a);
    // Put implementation
  virtual task T_out get();
    // Get implementation
endclass

8.25.3 Type Access

Parameters, constants and typedefs within a interface class are not inherited into the scope of descendent classes. All parameters and typedefs within an interface class are implicitly static and can be accessed through the class scope resolution operator :: (see 8.22). Nested classes (see 8.22) shall not be allowed in an interface class. An interface class shall not be nested within another interface class or within a class.

Some examples:

interface class fooIntf;
typedef enum {ONE, TWO, THREE} t1_t;
pure virtual function t1_t fooFunc();
endclass : fooIntf

class fooClass implements fooIntf;
t1_t t1_i; // error, t1_t is not inherited from fooIntf
virtual function fooIntf::t1_t fooFunc(); // correct. The scoping operator :: is used to access type t1_t
  return (fooIntf::ONE);
endfunction : fooFunc
endclass : fooClass

interface class interfaceClassA #(type T = logic)
typedef T[1:0] T2;
pure virtual function T2 foo();
endclass : interfaceClass

interface class interfaceClassB #(type T = int) extends interfaceClassA #(T);
endclass : interfaceClassB

interface class interfaceClassC #(type T = int) extends interfaceClassA #(T);
  pure virtual function T1 t1_bar(); // illegal, type T1 is not inherited into the scope of interfaceClassC
endclass : interfaceClassC

class derivedClass2 implements interfaceClassB #[bit];
  virtual function interfaceClassA::T2 foo(); // illegal, the return type is logic[1:0]; the inherited prototype return type is bit[1:0]
    // implement foo
endclass : derivedClass2

class derivedClass3 implements interfaceClassB #[bit];
virtual function interfaceClassA::(bit):T2 foo(); // legal, proper parameterization makes the types agree
  // implement foo
endclass : derivedClass3

class derivedClass4 implements interfaceClassB #[bit];
  virtual function bit[1:0] foo(); // legal, the return type bit[1:0] agrees with the inherited prototype
  // implement bar
endclass : derivedClass4

8.25.3.1 Type Usage Restrictions

There is a restriction placed on what type is passed along to the implementation of an Interface Class. An interface class type shall not be used in the parameterization of classes that implement them. The implemented class type shall be known at the point of reference to be an interface class. A class shall not implement a parameter, even if it resolves to an interface class. The following examples demonstrate this restriction and are illegal:

```verilog
class Fifo #(type T = PutImp) implements T;
virtual class Fifo #(type T = PutImp) implements T
interface class Fifo #(type T = PutImp) implements T;
```

Forward typedef of interface classes shall not be allowed. This means that all interface class definitions must be declared before the implementation of the interface class is reached.

```verilog
typedef class interfaceClassA; // Illegal forward typedef of an interface class
class derivedClass4 implements interfaceClassA #(bit);
  virtual function bit[1:0] foo();
  // implement bar
endclass : derivedClass4
```

// This interface class declaration must be declared before derivedClass4
```verilog
interface class interfaceClassA #(type T1 = logic)
typedef T1[1:0] T2;
pure virtual function T2 foo();
endclass : interfaceClass
```

8.25.4 Casting and Object reference assignment

There are a handful of relationships that must be clearly defined in order for interface classes to work properly with SystemVerilog. In order to maintain the OOP and polymorphic semantics, it shall be legal to assign an interface class handle to a child object that implements it.

```verilog
PutImp #() put_ref;
Fifo#() fifo_obj = new;
put_ref = fifo_obj;
```

It shall also be possible to have multiple references of an interface class and use them to cast from one to another.

```verilog
PutImp #() put_ref;
GetImp #() get_ref;
Fifo#() fifo_obj = new;
put_ref = fifo_obj;
$cast(get_ref, put_ref);
```
It shall also be legal to cast implemented objects onto their prototype interface handles

```plaintext
$cast(fifo_obj, put_ref); // This is legal
$cast(put_Ref, fifo_obj); // Legal, but casting is not required
```

Like abstract classes, a variable of an interface class type shall not be instantiated.

```plaintext
put_ref = new(); // This is illlegal
```

﴾shall be an error to cast from a source handle that is null. (See section 8.15 Casting)﴿

### 8.25.5 Name Scoping Conflicts and Resolution

When any class implements an interface class, the names and prototypes of that interface classes methods are declared as method names within the implementing class. They may be accessed in all ways that it is legal to access other names in that scope. Any subsequent definition, or duplicate declaration of these names must have exactly the same prototype (except that a definition eventually replaces the "pure" modifier). It is an error to declare, define, or acquire (from a super class or from another interface class) a method with the same name that does not exactly match the method prototype as declared in the interface class being implemented. The one exception to this exact matching rule allows derived classes to use a matching return type and not match the name exactly (see 8.19)

Let's take the following examples:

```plaintext
interface class interfaceBase;
  pure virtual function bit foo();
endclass

interface class interfaceExt extends interfaceBase;
  virtual function bit bar();
endclass

virtual class A implements interfaceBase;
  virtual function bit foo();
endclass

class derivedClass extends A implements interfaceExt;
  virtual function bit foo();
    return (0);
  endfunction
  virtual function bit bar();
    return (0);
  endfunction
endclass
```

In the above example, interfaceBase has a method named foo which is not implemented. interfaceExt is an extension of interfaceBase and adds bar as another unimplemented method. Class A then implements interfaceBase but is virtual so no implementation is required. Finally derivedClass extends A and implements interfaceExt forcing it to provide full definitions to both foo and bar. But notice the name collisions seen by derivedClass. It sees foo from the A class extension and the foo from the implements of interfaceExt. Because the prototype of the foo method is identical through both the extends and implements paths, derivedClass will only need to create the full definition of the prototyped method named foo.

The same example now shows how a collision can occur that results in an error:

```plaintext
interface class interfaceBaseA;
```
In this case, `foo` is defined in both `interfaceBaseA` and `interfaceBaseB` but with different result types, bit and int respectively. Because they are not identical prototypes, an error will be issued. The same will occur if a prototype with the same name is defined twice with different parameter values:

```verilog
interface class PutImp#(type T = logic, WIDTH = 1);
    pure virtual task void put(T [WIDTH-1:0] a);
endclass

interface class PutGetImp#(type T = logic);
    pure virtual task void put(T a);
    pure virtual task T get();
endclass

class Fifo#(type T = logic, WIDTH = 1, DEPTH = 1) implements PutImp(T, WIDTH),
    PutGetImp(T);
    T [DEPTH-1:0] myFifo;
    virtual task void put(T a);
    return (0);
endfunction
    virtual task T get();
    return (0);
endfunction
endclass
```

This example has the `PutImp` class with 2 parameters, `T` and `WIDTH`, feeding the `put` method within it. The `PutGetImp` class also has a definition of `put` but only has one parameter, `T`. When the `Fifo` class implements both of these interface classes it will recognize the parameter difference between the two versions of `put` and will issue an error.

A class may implement its interface class contracts by either defining or inheriting (or both) a good implementation of each unique pure method it promises to provide. The class which defines the method need not be the one to claim that it implements an interface class method. Here is an example:

```verilog
class baseClass;
    virtual function bit foo();
    return (1);
endfunction
endclass

interface class interfaceClass;
    pure virtual function bit foo();
    pure virtual function bit bar();
endclass

class derivedClass extends baseClass implements interfaceClass;
    // the “contract” to implement foo is fulfilled
    // by the inheritance from baseClass, even though baseClass
```
virtual function function bit bar();
    return (0);
endfunction
endclass

In this example baseClass fully defines foo, but foo is also prototyped in interfaceClass. derivedClass then extends baseClass and implements interfaceClass but because the prototype of foo in interfaceClass matches the prototype of the full definition of foo in baseClass, derivedClass uses foo defined in baseClass.

8.25.6 Partial implementation

It is possible to create classes that are not fully defined and which take advantage of interface classes through the use of virtual classes (see 8.20 Abstract Classes and pure virtual methods). Because virtual classes do not have to fully define their implementation, they are free to partially define their methods. The following is an example of a partially implemented virtual class.

class baseClass;
    virtual function bit foo();
        return (1);
    endfunction
endclass

interface class interfaceClass;
    pure virtual function bit bar();
endclass

virtual class derivedClass extends baseClass implements interfaceClass;
endclass

In this case derivedClass is virtual. It both extends baseClass which contains a fully defined method foo and implements interfaceClass which has only a prototype of bar.

8.25.7 Method default argument values

Method declarations within Interface Classes may have default argument values. The default expression is evaluated in the scope containing the subroutine declaration each time a call using the default is made. The actual value of the constant shall be the same for all the classes that implement the method. See section 13.5.3 (Default argument values) for more information.

8.25.8 Constraint Blocks and Cover Groups

Constraint blocks and cover groups shall not be allowed in interface classes.