

The "let" expression - motivation

Properties and sequences serve as templates for concurrent assertions, but there is no such construct for immediate assertions. Therefore, it is not possible to create packages/libraries of assertion templates for both concurrent and immediate assertions. The proposed **let** construct thus provides such a template as one of its possible uses. For example,

```
// in a package
let at_least_two(sig, rst = 1'b0) = rst || ($countones(sig) >= 2);

// in a design
reg [15:0] sig1; reg [3:0] sig2;
always_comb begin
    q1: assert (at_least_two(sig1));
    q2: assert (at_least_two(~sig2));
end
```

Note that in this case the **let** expression cannot be substituted by a function since formal arguments of a function need to have a specific type. There would have to be different function definitions for different argument widths. The **let** expression does not require that the type of the arguments be specified, it is inferred from the context.

The only alternative to **let** is a compiler directive. However, the **let** construct is safer because it has a local scope, while the scope of compiler directives is global. **let** is also more flexible, since like sequences and properties it allows defining default argument values and argument passing both by order and by name. Including **let** expressions into packages (See Clause 25) is a natural way to implement a well-structured customization for assertions. For example, the following **let** construct

```
package pex_gen9_common_expressions;
    let VALID_ARB(req, vld, arb_override) = (|(req&vld) || arb_override));
    ...
endpackage

module my_checker;
    import pex_gen9_common_expressions::*;
    logic a, b;
    wire [1:0] request;
    wire [1:0] valid;
    reg arb_out, ovr;
    ...
    prop: assert property (@(posedge clk)
        request |-> VALID_ARB(request, valid, ovr));
    ...
endmodule
```

is roughly equivalent to the following directive:

```
`define at_least_two(sig, rst = 1'b0) rst || ($countones(sig) >= 2)
```

Another use for **let** is in modeling for assertions instead of introducing additional nets. Usually RTL level code is used for this purpose. For example,

```
wire type(a + b) c; assign c = a + b;
...
assert property (@(posedge clk) cond |> c < d);
```

A synthesis tool may treat the auxiliary nets such as **c** as real ones and synthesize them into silicon. This may not be desirable. Workarounds using ``ifdef` are needed for this purpose. Also, the exact type must be

specified or the **type** operator must be used for this purpose, but it becomes awkward with long signal or expression names. Moreover, the type operator represents the self-determined width of the result while the user needs the context-determined width. Thus, if both *a* and *b* are of bit width two then *c* will also have the bit width two, and the result will sometimes be truncated. Using the **let** expression makes writing more elegant and safe:

```
let c = a + b; // let declaration
...
assert property (@(posedge clk) cond | => c < d);
```

The construct thus facilitates the creation of a modeling layer for SystemVerilog assertions and assertion-based checkers.

As a byproduct of this proposal adding a generate block to the list of the scopes where a sequence and a property may be declared.

14.3 Clocking block declaration

REPLACE in *Syntax 14-1*

```
clocking_item ::=
    default default_skew ;
    | clocking_direction list_of_clocking_decl_assign ;
    | { attribute_instance } concurrent_assertion_item_declaration
```

WITH

```
clocking_item ::=
    default default_skew ;
    | clocking_direction list_of_clocking_decl_assign ;
    | { attribute_instance } concurrent_assertion_item_declaration
```

16.6 Let construct (Note to the editor: Please shift clause numbering accordingly.)

```
assertion_item_declaration ::= // from A.2.10
    ...
    | let_declaration
let_declaration ::=
    let let_identifier ( ([let_port_list] ) ) = expression;
let_identifier ::=
    identifier
let_port_list ::=
    let_port_item { , let_port_item }
let_port_item ::=
    { attribute_instance } let_formal_type identifier [= expression]
let_formal_type ::=
    data_type_or_implicit
    | untyped This production depends on Mantis 1601
let_instance ::=
    let_identifier ( ([let_list_of_arguments] ) )
```

```

let_list_of_arguments ::=
    [let_actual_arg] {, [let_actual_arg] } {, . identifier ([let_actual_arg]) }
    | . identifier ([let_actual_arg] ) {, . identifier ([let_actual_arg] ) }
let_actual_arg ::=
    expression
expression ::=
    ...
    | let_instance40

```

//from A.8.3

Note to the editor: insert the first available reference number. Currently it is 40

40) *let_instance* may be used only in the expressions which are part of sequence or property expressions, **let** declarations or immediate assertions

Note to the editor: insert the first available reference number. Currently it is 40.

Syntax 16-2—Declaring *let* syntax

Note to the editor: Shift the numeration of the subsequent tables accordingly.

A **let** declaration defines a template expression (a **let** expression), customized by its ports. A **let** declaration may be instantiated in other **let** declarations, sequences, properties, concurrent and immediate assertions. A **let** declaration shall not be instantiated in the action block of a verification statement (see Clause 16.14).

let declarations can be used for customization and can replace the compiler directives in many cases. The **let** construct is safer because it has a local scope, while the scope of compiler directives is global. **let** is also more flexible, since like sequences and properties it allows defining default argument values, and argument binding by name as well as by position. Furthermore, including **let** statements into packages (see Clause 25) is a natural way to implement a well-structured customization for assertions. For example,

```

package pex_gen9_common_expressions;
    let VALID_ARB(req, vld, arb_override) = (! (req&vld) || arb_override);
    ...
endpackage

module my_checker;
    import pex_gen9_common_expressions::*;
    logic a, b;
    wire [1:0] request;
    wire [1:0] valid;
    reg arb_out, ovr;
    ...
    prop: assert property (@(posedge clk)
        request |-> VALID_ARB(request, valid, ovr);
    ...
endmodule

```

Similar to the way that properties and sequences serve as templates for concurrent assertions, the **let** construct can serve this purpose for immediate assertions. For example,

```

let at_least_two(sig, rst = 1'b0) = rst || ($countones(sig) >= 2);
reg [15:0] sig1; reg [3:0] sig2;
always_comb begin
    q1: assert (at_least_two(sig1));

```

```
    q2: assert (at_least_two(~sig2));  
end
```

In this case the `let` instantiation cannot be replaced by a function call since formal arguments of a function need to have a specific type. Thus there would have to be different function definitions for different argument widths. Another alternative to `let` is a compiler directive, but it has a global scope, as described earlier:

```
`define at_least_two(sig, rst) rst || ($countones(sig) >= 2)
```

Another intended use of `let` is in modeling for assertions to provide shortcuts for identifiers or subexpressions. For example,

```
let c = a + b;  
...  
assert property (@(posedge clk) cond | => c < d);
```

A `let` declaration may be instantiated in expressions of immediate assertions and match items of sequences, properties, verification statements, and other `let` declarations. A `let` declaration shall not be instantiated in the action block of a verification statement. The formal arguments may optionally be typed, following a similar syntactic form as for arguments to sequences and properties (see Clause 16.7.1) except that the argument types `sequence` and `property` cannot be used here (note to the editor: This refers to the updated clause from Mantis #1601). The argument types are restricted the same way as in boolean expressions used in assertions (see Clause 16.5.1). The formal arguments may have optional default values.

Like in sequences and properties, variables used in a `let` expression that are not formal arguments to the expression are resolved according to the scoping rules from the scope in which the `let` is declared.

In the scope of declaration, `let` expressions must be defined before used. No cross-module references to `let` declarations are allowed.

The `let` expression is substituted in the place of instantiation without any partial evaluation of the expression. When a `let` expression is instantiated, actual arguments can be passed to the expression. The `let` expression gets expanded with the actual arguments by replacing the formal arguments with the actual arguments. Semantic checks are performed to ensure that the expanded expression with the actual arguments is legal. The substituted expression is enclosed in parentheses (...) so as to preserve the priority of evaluation of the `let` expression. Recursive `let` instantiations are not permitted.

`let` expressions may contain sampled value function calls (`$rose`, `$fell`, `$past`, `$stable`, `$changed`). Their clock, if not explicitly specified, is inferred in the instantiation context in the same way as if the functions were used directly in sequences, properties or verification statements.

A `let` can be declared in any of the following:

- A module
- An interface
- A program
- A clocking block
- A package
- A compilation-unit scope
- A generate block

Examples:

- a) `let` with arguments and without arguments.

```
module m;
```

```

logic clk, a, b;
logic p, q, r;

// let with formal arguments and default value on y
let eq(x, y = b) = x == y;

// without parameters, binds to a, b above
let tmp = a && b;

...

a1: assert property (@(posedge clk) eq(p,q));
always_comb begin
    a2: assert (eq(r)); // use default for y
    a3: assert (tmp);
end

endmodule : m

```

The effective code after expanding **let** instances:

```

module m;
    bit clk, a, b;
    logic p, q, r;
    // let eq(x, y = b) = x == y;
    // let tmp = a && b;
    ...
    a1: assert property (@(posedge clk) m.p == m.q);
    always_comb begin
        a2: assert ((m.r == m.b)); // use default for y
        a3: assert ((m.a && m.b));
    end
endmodule : m

```

b) Declarative context binding of **let** arguments

```

module top;
    bit x = 1'b1;
    bit a;
    let y = x;
    ...
    always_comb begin
        // y binds to preceding definition of x
        // in the declarative context of let
        bit x = 1'b0;
        a1: assert (a || y);
    end
endmodule : top

```

The effective code after expanding **let** instances:

```

module top;
    bit x = 1'b1;
    bit a;
    // let y = x;
    ...
    always_comb begin
        // y binds to preceding definition of x
        // in the declarative context of let
        bit x = 1'b0;
        a1: assert (a || (top.x));
    end
endmodule : top

```

c) Sequences (and properties) with **let** in structural context

```
module top;
  parameter int p = 0;
  logic a, b;
  let x = a || b;
  sequence s;
    x ##1 b;
  endsequence : s
  ...
endmodule : top
```

The effective code after expanding **let** instances:

```
module top;
  parameter int p = 0;
  logic a, b;
  // let x = a || b;
  sequence s;
    (top.a || top.b) ##1 b;
  endsequence : s
  ...
endmodule : top
```

d) **let** declared in a **generate** block

```
module m(...);
  bit clk, a, b;
  bit [2:0] c;
  ...
  for (genvar i = 0; i < 3; i++) begin : L0
    if (i != 1) begin : L1
      let my_let(x) = !x || b && c[i];
      a1: assert property (@(posedge clk) my_let(a)); //OK
    end : L1
  end : L0
  a2: assert property (@(posedge clk) L0[0].L1.my_let(a)); // Illegal
endmodule : m
```

Assertion a1 after expansion becomes two assertions L0[0].L1.a1 and L0[2].L1.a1, the first of them being

```
assert property (@(posedge clk) (!m.a || m.b && m.c[0]));
```

and the second one being

```
assert property (@(posedge clk) (!m.a || m.b && m.c[2]));
```

Assertion a2 is illegal since it references the **let** expression hierarchically, while cross-module references to **let** expressions are not allowed.

e) **let** with typed arguments

```
module m(...);
  logic [15:0] a, b;
  typedef bit [15:0] bits;
  ...
  let ones_match(bits x, y) = x == y;
  always_comb
```

```

        al:assert(ones_match (a, b));
    endmodule : m

```

In this example assertion the `let` expression `ones_match` checks that both arguments have bits set to 1 at the same position. Because of the explicit specification of the formal argument types in the `let` declaration, all argument bits having unknown logic value or a high-impedance value become 0, and therefore the comparison captures the match of the bits set to 1. The effective code after expanding `let` instances:

```

module m(...);
    logic [15:0] a, b;
    typedef bit [15:0] bits;
    ...
    // let ones_match (bits x, y) = x == y;
    always_comb
        al:assert(bits'(a) == bits'(b));
endmodule : m

```

Note to the editor: the numeration of clauses and tables in the rest of this document is before shifting.

16.7 Declaring sequences

REPLACE

A **sequence** can be declared in

- A module
- An interface
- A program
- A **clocking** block
- A package
- A compilation-unit scope

WITH

A **sequence** can be declared in

- A module
- An interface
- A program
- A **clocking** block
- A package
- A compilation-unit scope
- A **generate block**

REPLACE in Syntax 16-4—Declaring sequence syntax

```

concurrent_assertion_item_declaration ::=
    ...
    | sequence_declaration

```

WITH

```

concurrent_assertion_item_declaration ::=
    ...

```

| sequence_declaration

...

16.12 Declaring properties

REPLACE

A **property** can be declared in any of the following:

- A module
- An interface
- A program
- A **clocking** block
- A package
- A compilation-unit scope

WITH

A **property** can be declared in any of the following:

- A module
- An interface
- A program
- A **clocking** block
- A package
- A compilation-unit scope
- A **generate block**

REPLACE in *Syntax 16-14—Property construct syntax*

```
concurrent_assertion_item_declaration ::=                               // from A.2.10
    property_declaration
    ...
```

WITH

```
concurrent_assertion_item_declaration ::=                             // from A.2.10
    property_declaration
    ...
```

25.2 Package declarations

REPLACE in *Syntax 25-1—Package declaration syntax*

```
package_or_generate_item_declaration ::=
    net_declaration
    | data_declaration
    | task_declaration
    | function_declaration
    | dpi_import_export
    | extern_constraint_declaration
    | class_declaration
    | class_constructor_declaration
```

```

| parameter_declaration ;
| local_parameter_declaration
| covergroup_declaration
| overload_declaration
| concurrent_assertion_item_declaration
|
;
```

WITH

```

package_or_generate_item_declaration ::=
net_declaration
| data_declaration
| task_declaration
| function_declaration
| dpi_import_export
| extern_constraint_declaration
| class_declaration
| class_constructor_declaration
| parameter_declaration ;
| local_parameter_declaration
| covergroup_declaration
| overload_declaration
| concurrent_assertion_item_declaration
|
;
```

A.1.10 Package items

REPLACE

```

package_or_generate_item_declaration ::=
net_declaration
| data_declaration
| task_declaration
| function_declaration
| dpi_import_export
| extern_constraint_declaration
| class_declaration
| class_constructor_declaration
| parameter_declaration ;
| local_parameter_declaration
| covergroup_declaration
| overload_declaration
| concurrent_assertion_item_declaration
|
;
```

WITH

```

package_or_generate_item_declaration ::=
net_declaration
| data_declaration
| task_declaration
| function_declaration
| dpi_import_export
| extern_constraint_declaration
| class_declaration
| class_constructor_declaration
```

```

| parameter_declaration ;
| local_parameter_declaration
| covergroup_declaration
| overload_declaration
| concurrent_assertion_item_declaration
| ;

```

A.2.10 Package items

REPLACE

```

concurrent_assertion_item_declaration ::=
    property_declaration
    | sequence_declaration

```

WITH

```

concurrent_assertion_item_declaration ::=
    property_declaration
    | sequence_declaration
    | let_declaration

```

ADD at the end

```

let_declaration ::=
    let let_identifier ( ([let_port_list]) ) = expression;

```

```

let_identifier ::=
    identifier

```

```

let_port_list ::=
    let_port_item { , let_port_item }

```

```

let_port_item ::=
    { attribute_instance } let_formal_type identifier [= expression]

```

```

let_formal_type ::=
    data_type_or_implicit
    | untyped      This production depends on Mantis 1601

```

```

let_instance ::=
    let_identifier ( ([let_list_of_arguments]) )

```

```

let_list_of_arguments ::=
    [let_actual_arg] { , [let_actual_arg] } { , . identifier ([let_actual_arg]) }
    | . identifier ([let_actual_arg]) { , . identifier ([let_actual_arg]) }

```

```
let_actual_arg ::=  
    expression
```

A.6.11 Clocking block

REPLACE

```
clocking_item ::=  
    default default_skew ;  
    | clocking_direction list_of_clocking_decl_assign ;  
    | { attribute_instance } concurrent_assertion_item_declaration
```

WITH

```
clocking_item ::=  
    default default_skew ;  
    | clocking_direction list_of_clocking_decl_assign ;  
    | { attribute_instance } concurrent_assertion_item_declaration
```

A.2.10 Assertion declarations

ADD at the end

A.2.10 Expressions

REPLACE

```
expression ::=  
    primary  
    | unary_operator { attribute_instance } primary  
    | inc_or_dec_expression  
    | ( operator_assignment )  
    | expression binary_operator { attribute_instance } expression  
    | conditional_expression  
    | inside_expression  
    | tagged_union_expression
```

WITH

```
expression ::=  
    primary  
    | unary_operator { attribute_instance } primary  
    | inc_or_dec_expression  
    | ( operator_assignment )  
    | expression binary_operator { attribute_instance } expression  
    | conditional_expression  
    | inside_expression  
    | tagged_union_expression  
    | let\_instance40
```

Note to the editor: insert the first available reference number. Currently it is 40.

A.10 Footnotes (normative)

ADD to the end

40) *let_instance* may be used only in the expressions which are part of sequence or property expressions, let declarations or immediate assertions

Note to the editor: insert the first available reference number. Currently it is 40.

Add to Table B1—Reserved keywords

let

Annex M

Modify M.2

from

```
/* property decl, spec */
#define vpiPropertyDecl 655
#define vpiPropertySpec 656
#define vpiPropertyExpr 657
#define vpiMulticlockSequenceExpr 658
#define vpiClockedSeq 659
#define vpiPropertyInst 660
#define vpiSequenceDecl 661
#define vpiActualArgExpr 663
#define vpiSequenceInst 664
#define vpiImmediateAssert 665
#define vpiReturn 666
```

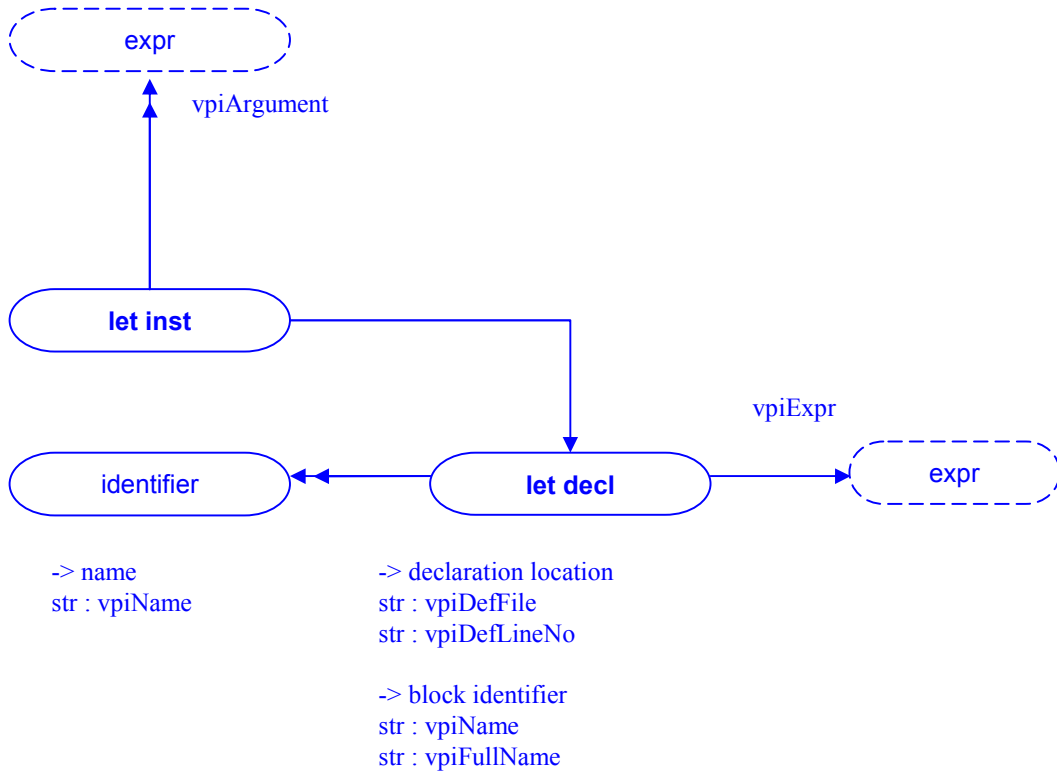
to

```
/* property decl, spec */
#define vpiPropertyDecl 655
#define vpiPropertySpec 656
#define vpiPropertyExpr 657
#define vpiMulticlockSequenceExpr 658
#define vpiClockedSeq 659
#define vpiPropertyInst 660
#define vpiSequenceDecl 661
#define vpiActualArgExpr 663
#define vpiSequenceInst 664
#define vpiImmediateAssert 665
#define vpiLetDecl Editor to fill
#define vpiLetInst Editor to fill
#define vpiReturn 666
```

Add to Clause 36

36.76 let declaration

Note to Editor: It is best to put this after 36.48 Multi-clock sequence, and renumber all that follows.



Details :

The vpiArgument iterator shall return the let instance arguments in the order that the formals for the let are declared, so that the correspondence between each argument and its respective formal can be made. If a formal has a default value, that value shall appear as the argument should the instantiation not provide a value for that argument.

36.50 Expressions

CHANGE TO

