



## SUPERLOG<sup>®</sup> Design Assertion Subset

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## Change History

Version 1.6 of this document has been updated from version 1.5 to reflect the results of the Accellera SystemVerilog Assertions committee meetings held on

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## Introduction

An assertion is a statement that a property must be true. There are two kinds of assertions: concurrent assertions which state that the property must be always be true, e.g. throughout a simulation, and procedural assertions which are incorporated in procedural code and apply only for a limited time or under limited conditions.

There are various applications of assertions. They can be included in the design, to document the assumptions made by the designer and to facilitate "white box" testing. They can be outside the design, either in a testbench to check the response of the design to the stimulus, or to control a tool such as a stimulus generator or a -model checker.

Concurrent assertions can be coded as modules in a library, but this limits the complexity of the property that can be expressed easily. It is more difficult to code procedural assertions as a library of tasks in Verilog, because events cannot be arguments, each assertion may need static data, and tasks block.

## Procedural Assertions

SUPERLOG Design Assertions Subset provides ~~the following~~ four kinds of procedural assertions, which allow the user to test boolean expressions or sequences of boolean expressions, and perform some action based on whether the expression or sequence is true or false. *Immediate* assertions test the value of a boolean expression at the time the statement is executed, and may be used in always and initial blocks, tasks and functions. *Strobed* assertions schedule the evaluation of the expression to be delayed until the end of the current timeslice, to allow for glitches to settle. Strobed assertions may be used in initial and always blocks and tasks, but not in functions since functions must return immediately.

To test sequences of expressions, it is necessary to specify a *sampling clock* event on which to test each element of the sequence. Therefore, *Clocked Immediate* and *Clocked Strobed* assertions are added to allow progressive evaluation of sequences of expressions. Since these clocked assertions, by definition, take time, they cannot be used in functions. Clocked immediate assertions evaluate each expression in the sequence when the clock event triggers, and clocked strobed assertions evaluate each expression at the end of the timeslice at which the event triggers.

~~Immediate assertions~~

~~Strobed assertions~~

~~Clocked immediate assertions~~

~~Clocked strobed assertions~~

~~These are statements in an initial or always block, or in a task. Immediate assertions can also be in a function.~~

## Syntax

```
<proc_assertion> ::= [<identifier> ':' ] <immediate_assert>
                  | [<identifier> ':' ] <strobed_assert>
                  | [<identifier> ':' ] <clocked_immediate_assert>
                  | [<identifier> ':' ] <clocked_strobed_assert>

<immediate_assert> ::= 'assert' '(' <expression> ')'
                    <statement_or_null> //pass
                    ['else' <statement_or_null>] //fail

<strobed_assert> ::= 'assert_strobe' '(' <expression> ')'
                  <restricted_statement_or_null> //pass
                  ['else' <restricted_statement_or_null>] //fail

<clocked_immediate_assert> ::= 'assert' [<reset>]
                              '(' <formula_expr> ')' <step_control>
                              <statement_or_null> //pass
                              ['else' <statement_or_null>] //fail

<clocked_strobed_assert> ::= 'assert_strobe' [<reset>]
                             '(' <formula_expr> ')' <step_control>
                             <restricted_statement_or_null>
                             ['else' <restricted_statement_or_null>] //fail

<formula_expr> ::= <expr_sequence>
                 | <expr_sequence> 'triggers' <formula_expr>

<expr_sequence> ::= <expression>
                  | '[' <constant_expression> ']' // skip n steps
                  | <range> // skip m to n steps
                  | <expr_sequence> ';' <expr_sequence> // sequence
                  | <expr_sequence> '*' '[' <constant_expression> ']'
                    //repetition
                  | <expr_sequence> '*' <range> // bounded repetition
                  | '(' <expr_sequence> ') '

<range> ::= '[' <constant_expression> ':' <constant_expression> ']'

<step_control> ::= '@@' <name>
                 | '@@' '(' <event_expressions> ') '

<reset> ::= 'accept' '(' <expression> ')' // sync
           | 'accept' <event_control> // async
           | 'reject' '(' <expression> ')' // sync
           | 'reject' <event_control> // async
```

## Immediate Assertions

The immediate assert statement is a test of an expression performed when the statement is executed in the procedural code. The expression is treated as a condition like in an if statement:

```
[<identifier> ':'] assert (<expression>) [<pass_statement>] [else <fail_statement>]
```

The pass statement is executed if the assertion succeeds, i.e. the expression evaluates to 'true'. As with the 'if' statement, if the expression evaluates to 'X', 'Z' or '0', then the assertion fails. The pass statement may, for example, record the number of successes for a coverage log, but may be omitted altogether. If the pass statement is omitted, then no action is taken if the assert expression is true. The fail statement is executed if the assertion fails (i.e. the expression does not evaluate to a known, non-zero value) and can be omitted. The optional assertion label (identifier and colon) creates a notional named block around the assertion statement (or any other SystemVerilog statement) and can be displayed using the %m format code.

```
assert_foo : assert (foo) $display("%m passed"); else $display("%m failed");
```

If the **else** is omitted a default error message is written if the assertion fails. The contents of the message may be tool-specific. For simulation tools, it is recommended that the message include basic information about the assertion statement and when it failed to facilitate debug, such as

```
Run-time WarningError: myfile.slg:7 Assertion assert_foo failed at time 105
```

If the **else** is present, the default message is still printed when the assertion fails. Additional information may be displayed after the default message by including a system task to do so, such as \$fatal, \$error, \$warning, \$info. These system tasks indicate the severity, and can contain additional information in the format of \$display.

```
assert (foo); else $fatal("finish simulation because foo is false");
```

\$fatal displays a Run-time Fatal, which terminates the simulation with an error code. The first argument passed to \$fatal shall be consistent with the argument to \$finish.

\$error displays a Run-time Error, which cannot be turned off. This shall be the default.

\$warning displays a Run-time Warning, which can be turned off.

\$info: displays an information message, which can be turned off by "quiet mode".

**\$quiet: no arguments. Suppresses the printing of the default message.**

**If more than one of these system tasks is present in the else clause, then all of them will be executed.**

```
<assert_msg> ::= <fatal_func> | <msg_func>;

<fatal_func> ::= $fatal ('<finish_num>', ' // finish arg
                        [<format_string>]', '<expr>...')';

;

<msg_func> ::= <msg_name> (' [<format_string>]', '<expr>...')'';

<msg_name> ::= $error | $warning | $info
```

The display of the default message and additional messages of **specific-warning and info** types may be controlled by a command-line option.

Since the fail statement, like the pass statement, is any legal SystemVerilog procedural statement, it may also be used to signal a failure to another part of the testbench:

```
assert (myfunc(a,b)) count1 = count + 1; else ->event1;
assert (y == 0); else flag = 1;
```

The assert statement serves as guidance to non-simulation tools that the condition should be true. The second statement above is equivalent to:

```
if ( y!=0) begin flag = 1; end
```

## Strobed Assertions

If an immediate assertion is in code triggered by a timing control that happens at the same time as a blocking assignment to the data being tested, there is a risk of the wrong value being sampled. For example:

```
always @(posedge clock) a = a + 1; // blocking assignment
always @(posedge clock) begin
    ....
    assert (a < b);
end
```

This can be solved by using a strobed assertion, which waits in the background until the end of the time slot, like the \$strobe system task.

```
always @(posedge clock) begin
    ....
    cas:assert_strobe (a < b);
end
```

Strobed assertions can have pass or fail statements like immediate assertions. However, they are restricted to another assertion statement, a system task call, a statement preceded by a delay control or an event control, or sequential block containing them. This is because the statement happens after the assertion is evaluated, at the end of the time slot, and hence ~~they~~ must not create more events at that time slot or change values. Statements which cause additional events to occur at the current time shall be an error.

The example below illustrates the effect of blocking and non-blocking assignments on immediate and strobed assertions. The immediate assertions are like \$display statements and the strobed assertions are like \$strobe statements:

```
module test;
reg [3:0] a=0; c=0, d=0;
reg clk = 0;
wire b;

initial begin
    #10 clk = 1;
    forever #5 clk = !clk; // posedge clk at 10,20,30,40...
end

assign b = a+1;

always @(posedge clk) begin
    a1: assert(c<3); // fails at time 40
    c = c+1;
    a2: assert(c<3); // fails at time 30
    a <= a+1;
    a3: assert(a<3); // fails at time 40
    a4: assert(b<3); // fails at time 40
    a5: assert_strobe(a<3); // fails at time 30
    a6: assert_strobe(b<3); // fails at time 30
end
```

```

always @(a) begin // models transient behavior on comb. nets
    d = a+2; // spikes to 2 at 0, 3 at 10, 4 at 20
    assert(d<3); // fails at time 10
    d = d-1; // settles to 1 at 0, 2 at 10, 3 at 20
    assert(d<3); // fails at time 20
end

always @(d) assert_strobe (d<3); // fails at time 20

endmodule

```

## Sequential Assertions

In addition to assertions about single expressions, it is often useful to assert sequences of expressions over time. One way of doing this is to use nested immediate assertions:

```

always @(posedge clk or negedge rst)
  if(state == REQ)
    a7: assert(req1) // no semicolon
        @(posedge clk) assert (gnt) @(posedge clk) assert(!req1);

```

The above example verifies the sequence that, if state is equal to REQ, the req1 signal must be true immediately, then on the next posedge clk, gnt must be true and on the following posedge clk, req must be false. Note that the assertion statement itself is nonblocking, so the sequence in assertion a7 is equivalent to

```

always @(posedge clk or negedge rst)
  if(state == REQ)
    a8: assert(req1)
        process // Please see the SystemVerilog Spec, section 10,
            // for more details about the process statement
            @(posedge clk) assert (gnt) @(posedge clk) assert(!req1);

```

To simplify this complex nested assertion, a *sequential regular expression* is used in the assert statement. Sequential regular expressions require a *step control* event expression to specify the timing between evaluations of each element in the regular expression. Using a sequential regular expression, the assertion a8 could be rewritten as

```

always @(posedge clk or negedge rst)
  if(state == REQ)
    a9: assert(req1;gnt;!req1) @@(posedge clk); // note the @@ token to distinguish the step control
        // from the pass statement

```

A sequential regular expression is a semicolon-delimited list of expressions. The first expression in the list is evaluated immediately when the assert statement is executed. The other subsequent expressions are evaluated one at a time on successive occurrences of the step control event expression. In assertion a9 above, req1 is evaluated immediately when the assert statement is executed, just as for an immediate assertion.

The '@@' token is introduced to distinguish the step control from an ordinary event control at the start of the pass statement. Consider the following:

```

always @(posedge clock or negedge rst)
  if(state == REQ)
    assert (req1)
        @(posedge clk) // This is an event control in the pass statement
            $display("Hello at time %t", $time);

```

In this example, the "@(posedge clock)" in the pass statement causes the display action to occur on the next posedge clock after the assertion succeeds. Therefore, a new token is required to distinguish the assertion sequence step control from the pass statement.

Note that, since the first expression is evaluated immediately, assertion a9 above is equivalent to

```
always @(posedge clk or negedge rst)
  if(state == REQ)
    assert(req1)
    assert(1;gnt;!req1) @@(posedge clk);
```

The sequence notation "(1;<expression\_or\_sequence>)" is a convenient shorthand indicating that the <expression\_or\_sequence> is to be evaluated on the next occurrence of the step control event. This is because the expression '1' is evaluated immediately and is always true.

Sequential assertions using the **assert** keyword are called *clocked immediate* assertions, since the expressions are evaluated as with immediate assertions. Similarly, *clocked strobed* assertions may be written using the **assert\_strobe** keyword, in which each expression in the sequence is evaluated at the end of the timeslice in which the assertion is executed or in which the step control event occurs. The pass and fail statements of clocked strobed assertions have the same restrictions as strobed assertions.

Specifying an explicit step control for a sequence makes it possible to use clocked assertions in combinational always blocks:

```
always @(foo,bar)
  assert_strobe (a;b;c) @@(posedge clk); // look for a when foo or bar change,
  // then look for b on next posedge clk
```

Since it is common for combinational always blocks to be executed multiple times in a single timestep as the signals in the event trigger expression settle, it is common to use strobed assertions be used in combinational always blocks. Immediate assertions are commonly used in clocked always blocks.

### ***Clocked Immediate Assertions***

~~A sequence of expressions can be asserted by specifying a *sequential regular expression* in the assert statement, along with a step control to specify the timing between evaluations of the sequential regular expression:~~

```
——— always @(posedge clock or negedge reset)
——— assert (a; b; c) @@(posedge clock); // note the @@ token to distinguish the step control
——— // from the pass statement
```

~~A sequential regular expression is a semicolon delimited list of expressions. The first expression in the list is evaluated immediately when the assert statement is executed. The other subsequent expressions are evaluated one at a time on successive occurrences of the step control event expression. In the above example, the expression 'a' is evaluated immediately, just as for an immediate assertion, with 'b' being evaluated on the next posedge clock and 'c' evaluated on the posedge after that. The immediate clocked assertion expressions are evaluated among the first active events in the timestep in which the explicit event is triggered.~~

~~The '@@' token is introduced to distinguish the step control from an ordinary event control at the start of the pass statement. Consider the following:~~

```
——— always @(posedge clock)
——— assert (a)
——— @(posedge clock) // This is an event control in the pass statement
——— $display("Hello at time %t", $time);
```

~~In this example, the "@(posedge clock)" in the pass statement causes the display action to occur on the next posedge clock after the assertion succeeds. Therefore, a new token is required to distinguish the assertion sequence~~

step control from the pass statement. Specifying an explicit step control for a sequence makes it possible to use clocked assertions in combinational always blocks:

```
——— always @(foo,bar)
——— assert (a;b;c) @@@(posedge clk); // look for a when foo or bar change,
——— // then look for b on next posedge clk
```

The assertion

```
——— always @(posedge clk)
——— a1: assert (a;b;c) @@@(posedge clk);
```

is nearly equivalent to

```
——— always @(posedge clk)
——— a2: assert (a) @@@(posedge clk) // no semicolon
——— @@@(posedge clk) a3: assert (b;c) @@@(posedge clk);
```

The assertion **a1** and the combination of assertions **a2** and **a3** both evaluate the same sequence, but splitting the sequence between two assertions allows the user to execute different fail statements depending on where the assertion fails. For example, it may be deemed an error if 'a' is not true now, but only a warning if 'b' is not true on the next clock edge. Another way of expressing the same behavior is

```
——— always @(posedge clk)
——— a4: assert (a) // no semicolon
——— a5: assert (1;b;c) @@@(posedge clk); // b evaluated on next posedge clk, then c on next posedge clk
——— else $warning("b or c not found"); // else clause for a5
——— else $error("a not found immediately"); // else clause for a4
```

In this example, assertion **a5** evaluates the first expression in the sequence immediately. The value '1' is always true, so on the next posedge of clk, it will evaluate 'b'. This is a convenient notation for specifying a sequence that begins on the next clock.

Note that to avoid races, the variables read in clocked immediate assertions should be written by non-blocking assignments. Expressions in clocked strobed assertions are always sampled at the end of the timestep, so no race conditions should occur.

An assertion could be executed twice in the same timestep via a zero-delay loop or a combinational always block, for example. If a clocked immediate assertion is executed more than once at the same timestep, the first expression in the sequence will be re-evaluated. If a clocked strobed assertion is executed more than once at the same timestep, the first expression in the sequence will be evaluated once at the end of the timestep.

An assertion shall only spawn a single process to evaluate the next expression in the sequence at the next step control event. If the step control event occurs multiple times at the same timestep, then in a clocked immediate assertion the current expression in the sequence shall be re-evaluated. In a clocked strobed assertion, the current expression will still be evaluated only once at the end of the timestep. The next expression in the sequence shall not be evaluated until the step control occurs in a later timestep.

As mentioned above, the execution of a sequential assertion spawns a process that monitors each event in the sequence when the step control event occurs. If the sequential assert statement is executed again before the sequence spawned by the original execution has expired, then a new process shall be spawned that looks for the sequence starting at the current timestep. It is therefore possible to have multiple processes in-flight, each monitoring the same sequence but offset in time. It is possible for these multiple processes to be satisfied by the same sequential behavior, even though the processes are offset in time. In such a case, both processes will terminate at the same timestep, in which both sequences are satisfied.

~~A clocked assertion can be triggered twice at the same timestep, and the first expression in the sequence will be re-evaluated. This could happen in a zero-delay loop, but is not recommended practice.~~

~~A clocked assertion may also be re-triggered at the next timestep, before the sequence has expired. Any single assertion shall only be spawned once at a particular timestep.~~

~~If the step control event occurs multiple times at the same timestep the evaluation shall not progress through the sequence. Instead, the current expression in the sequence is re-evaluated. Consider:~~

Consider:

```
module top;
  reg clk = 0;
  reg a,b,c;

  initial begin
    #10 clk = 1;
    forever begin
      clk = 0;
      clk = 1; // 2 posedges clk at 10,20,30,40...
      #5 clk = 0;
      #5 clk = 1;
    end
  end

  always @(posedge clk)
    assert(a;b;c) @@(posedge clk); // 'a' is evaluated only once at 10, 'b' once at 20, 'c' once at 30
```

Note that the step control expression may be any valid event expression in SystemVerilog. The following assertions all use valid step control expressions:

```
bit clk;
event ev1;

always @(posedge clk or negedge reset) begin
  assert (a;b;c) @@(negedge clk); // sequence sampled on negedge clk
  assert (a;b;c) @@(clk); // sequence sampled on any edge of clk
  assert (a;b;c) @@(ev1); // sequence sampled when event ev1 fires
end
```

This flexibility also allows nested assertions to use different clocks:

```
always @(posedge clk) begin
  assert (a;b) @@(posedge clk) // on posedge clk
  assert (1;c;d) @@(negedge clk); // look for c and d on negedge clk
  assert (e;f) @@(posedge clk2)
  assert (1;g;h) @@(ev1);
end
```

### **Clocked Strobed Assertions**

~~Another way to avoid races is to use the same constructs with strobed assertions:~~

```
always @(posedge clock) begin
  ...
  assert_strobe (a; b; c) @@(posedge clk); // a is sampled at end of current time step, b at end of next
  clock time step
  ...
end
```

~~————end~~

This is also equivalent to

~~————assert\_strobe (a) // no semicolon  
————assert\_strobe (1;b; c) @@@(posedge clk); // pass statement, if first one OK~~

This is also equivalent to:

~~————assert\_strobe (a) // no semicolon  
———assert\_strobe (1;b) @@@(posedge clk) // pass statement, if first one OK  
———a6: assert\_strobe (1;c) @@@(posedge clk); // pass statement, if second one OK~~

Note that assertion a6 will not evaluate until the successful completion of (a;b), which is why only a single '1' is required to specify that c be evaluated on the next posedge clk.

Strobed assertions with explicit clocks can be used in always blocks that model combinational logic and therefore do not have a clock and are subject to glitches.

Even if a clocked strobed assertion is triggered twice at the same clock cycle, it only executes once at the end of the time step. A clocked strobed assertion may be re-triggered at the next clock cycle, before the sequence has expired.

## More Expression Sequences

A number of steps can be skipped either by writing expressions which are always true:

**assert** (a;1;1;c) @@@(posedge clk); // two steps between a and c

or by using the notation [n] to count the number of steps:

**assert** (a;[2];c) @@@(posedge clk); // two steps between a and c  
**assert** (a;[1];[1];c) @@@(posedge clk); // two steps between a and c

Note that in [n], the n must be a **non-negative** literal or a constant expression. [0] has no effect. The number of steps to be skipped may also be expressed using [min:max], where the minimum number of steps **may** must be greater than or equal to zero. Both min and max must be a literal or constant expression.-

**assert** (a;[0:10];b) @@@(posedge clk); // b occurs between the next and 11<sup>th</sup> clock edges, inclusive

If an expression must be repeated a defined number of times, this can be expressed with a trailing \*[n]. If it can be repeated a minimum or maximum number of times, this can be expressed with a trailing \*[min : max]. These repetition counts must also be literals or constant expressions:

**assert** (a; b)\*[5]; // a;b;a;b;a;b;a;b;a;b  
**assert** (a\*[0:3];b;c); // (a;b) or (a;b;a;b) ~~**assert** (a; b)\*[1:2]; // (a;b) or (a;b;a;b)~~

Note that (a\*[0:3];b;c) is equivalent to (b;c) or (a;b;c) or (a;a;b;c) or (a;a;a;b;c). This means that a sequence a;ab;a;b;c; will pass. The expression sequence is not equivalent to ((a && !b)\* [0:3];b;c), which would fail the same sequence.

The rules for specifying repeat counts are summarized as:

- Each form of repeat count specifies a minimum and maximum number of occurrences:
  - expr\*[n:m] n is the minimum, m is the maximum
  - expr\*[n] same as expr\*[n:n]
  - [n] same as 1\*[n:n]
- The sum of the minimum repeat counts for all terms in a sequence must be greater than 0.
- The sequence as a whole cannot be empty.
- The last term in a sequence may not have a min:max range of repetition. If it does, it shall be an error.

## Antecedent and Consequent

An assertion is usually only applicable under a certain condition, called an antecedent. One way of specifying a condition is to use procedural code **if** or **case**. For example:

```
——— if (a) assert (1;b;c) @@(posedge clk); // only apply assertion if a is true
```

Another way is to use the **triggers** keyword in the assertion itself. The following is equivalent to the example above:

```
——— assert (a triggers (1;b;c)) @@(posedge clk); // only check b;c if a is true
```

If a is false when the assertion is applied, the assertion is aborted. It neither passes nor fails. If the antecedent a is true, the rest of the assertion, called the consequent, is applied.

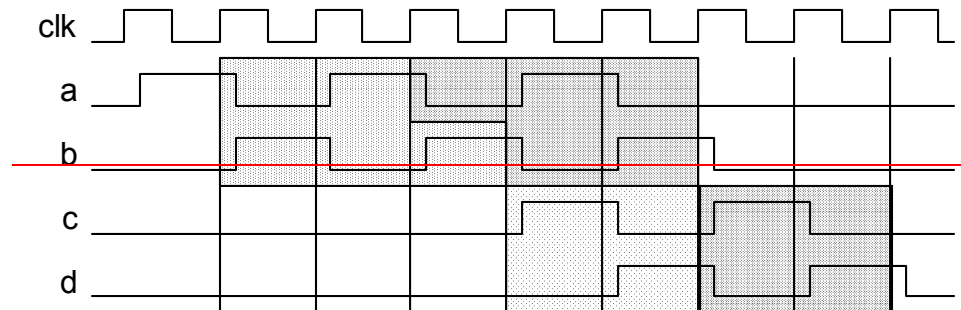
The benefit of using the second form is that the antecedent can be a sequence. For example:

```
——— assert ((a;b) triggers (1;c)) @@(posedge clk); // only check c if a is true then b is true
```

This technique can conveniently be used for assertions that are not embedded in procedural code, but are stand-alone concurrent statements

```
——— always @(posedge clock)
——— assert_strobe ((a;b;a;b) triggers (c;d)) @@(posedge clk);
```

Note that the antecedent (a;b;a;b) is retriggerable, so that the input sequence ababab requires two instances of the consequent (c;d) i.e. cded.



## Resetting Assertions

A named assertion can be disabled like any other named SUPERLOG block. If this is done before the expression sequence has finished, it means that neither the pass statement nor the fail statement is executed.

```
disable cas;
```

Note that if the disable is applied at the same simulation time step as the last clock step of a sequence, there is a race in the case of an immediate assertion, but a strobed assertion is always disabled.

An alternative way to reset the assertions is to use the **accept** or **reject** keywords. Acceptance means that the assertion succeeds, and the pass statement is executed. Rejection means that the assertion fails and the fail statement is executed. These keywords can be used either with a Boolean expression for synchronous reset, or with an event control for asynchronous reset. The following examples use an explicit clock:

```
a7: assert reject (!busy) (a;b;c) @@(posedge clk); // busy must be 1
assert accept (sync_reset) (a;b;c) @@(posedge clk);
assert reject @(negedge busy) (a;b;c) @@(posedge clk);
assert accept @(posedge async_reset) (a;b;c) @@(posedge clk);
assert_strobe reject (!busy) (a;b;c) @@(posedge clock); // busy must be 1
assert_strobe accept (sync_reset) (a;b;c) @@(posedge clock);
```

```

assert_strobe reject @(negedge busy) (a;b;c) @@(posedge clock);
assert_strobe accept @(posedge async_reset) (a;b;c) @@(posedge clock);

```

The effect of a synchronous **reject** expression is to logically and the inverse of the reject expression with each element in the sequence. Therefore, assertion a7 above is equivalent to

```

assert (a && busy; b && busy; c && busy) @@(posedge clk);

```

If an asynchronous reset occurs simultaneously with the last clock step, there is a race in the case of an immediate assertion and the reset action always happens in the case of a strobed assertion.

Other examples of synchronous accept/reject are shown here:

```

always @(posedge clk)
  if (fifoFsmState == READ)
    assert accept(ack) (dataValid * [15]) @@(posedge clk);

```

Once the fifo goes to READ state, the dataValid must be held high for 15 cycles or when the ack goes high, whichever comes first.

```

always @(arbiterState or request)
  if (arbiterState == GRANT)
    assert reject(~request) ((arbiterState==GRANT)*[5]) @@(posedge clk);

```

Once the arbiter goes to GRANT state, it must stay in that state for at least next 5 cycles during which the request signal cannot be lowered. If request goes low, the assertion fails.

### ***Issue For Discussion:***

*If an assertion is named, could it be accepted or rejected explicitly from a different block, similar to **disable**?*

*Example:*

```

always @(posedge clk)
  if(state == READ)
    a8: assert (!as;!ds;rdy*[0:10];!rdy;rdy) @@(posedge clk);

```

```

always @(posedge clk)
  case(arbState)
    ...
    UNFAIR: reject a8;
  endcase

```

## **Controlling Assertions**

System functions are provided to limit assertion checking to part of the design and part of the simulation time.

```

$assertoff (hierarchical_names) // name of module instance or individual assertion
$asserton (hierarchical_names) // name of module instance or individual assertion

```

Assertions are on by default until turned off. If the list of names is empty, it is taken to refer to all assertions.

The effect of turning assertions off is to stop the check and both the pass and fail statements. Assertions already started are not affected. They can be turned off individually by **disable** statements.

### **Controlling the Steps**

To control the stepping of always blocks, SystemVerilog has the following enhancement from Verilog:

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`@(posedge clk iff enable == 1)`

The event expression only triggers if the expression after the **iff** is 1, in this case when `enable == 1`. Note that such an expression is evaluated when `clk` changes not when `enable` changes. Also note that **iff** has precedence over **or**, so that the **iff** needs to be repeated for each **ored** event to which it applies.

Since this feature is a valid event expression, it can be used to control the stepping of an assertion sequence. In effect, this allows a "gated clock" to control the assertion without the user having to declare the gated clock explicitly. Because this could have significant impact on the ability of Formal Verification tools to evaluate the assertion successfully, it is recommended that this construct be used only for simulation.

## System Functions

Assertions are commonly used to evaluate certain specific characteristics of a design implementation, such as whether a particular signal is "one-hot." The following system functions are included to facilitate such common assertion functionality:

`$onehot (<expression>)// only one bit of <expression> is high`  
`$in (member, list) // member appears in the list`

### *Item for Discussion:*

*What other common actions are done in assertions that should have accompanying system functions?*

Concern: Must be implemented in a way that does not impact simulation performance.  
Orthogonal to the rest of the document.